

# Gray's tensor product of 2-categories

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There is a particularly simple locally-ordered 2-category  $\mathcal{Q}(T)$  obtained from each totally ordered set  $T$ . After describing this 2-category we shall produce a presentation which shows that, for  $T$  of cardinality  $N$ , it is Gray's free  $N$ -dimensional cube with commutative 3-dimensional faces [1].

Gray [2] proved that the 2-category with this presentation is locally ordered without describing an explicit model such as our  $\mathcal{Q}(T)$ . He made use of Artin's braid group [3]. Our approach uses the normal form for elements of the braid monoid with zero which we learned from Eilenberg [4]. (Eilenberg - Street [5] have shown that this normal form can be obtained from a "rewrite system".)

Let  $T$  be any totally ordered set. For any list  $a = a_1 a_2 \dots a_r$  of elements of  $T$ , let

$$\sigma(a) = \{(a_i, a_j) \mid i < j, a_i < a_j\}.$$

The relation between  $\sigma$  and concatenation of lists is

$$\sigma(ab) = \sigma(a) + \{(a_i, b_j) \mid a_i < b_j\} + \sigma(b). \quad (1)$$

The objects of the 2-category  $\mathcal{Q}(T)$  are finite subsets  $x$  of  $T$ . For  $y \subseteq x$ , the ordered set  $\mathcal{Q}(T)(x, y)$  consists of listings  $a$  of all the elements of  $x - y$  together with the order:

$$a \leq a' \quad \text{when} \quad \sigma(a) \subseteq \sigma(a').$$

The other homs of  $\mathcal{Q}(T)$  are empty. Composition

$$\mathcal{Q}(T)(x, y) \times \mathcal{Q}(T)(y, z) \longrightarrow \mathcal{Q}(T)(x, z)$$

is concatenation of lists, which is order preserving by equation (1).

Let  $\mathcal{G}(T)$  denote the graph whose vertices are finite subsets  $x$  of  $T$  and whose edges  $u: x \rightarrow y$  are elements  $u$  of  $x$  such that  $x = y + \{u\}$ . We can identify the free category on  $\mathcal{G}(T)$  with the underlying category of  $\mathcal{Q}(T)$ .

A graph together with some 2-cells between  $c$ terminal paths in the graph was called a computed by Street [ ].

Enrich  $\mathcal{G}(T)$  with the structure of computed by giving it a 2-cell

$$\begin{array}{ccc} x & \xrightarrow{u} & x - \{u\} \\ v \downarrow & \langle u, v \rangle \rightrightarrows & \downarrow v \\ x - \{v\} & \xrightarrow{u} & x - \{u, v\} \end{array} \quad u < v \quad (2)$$

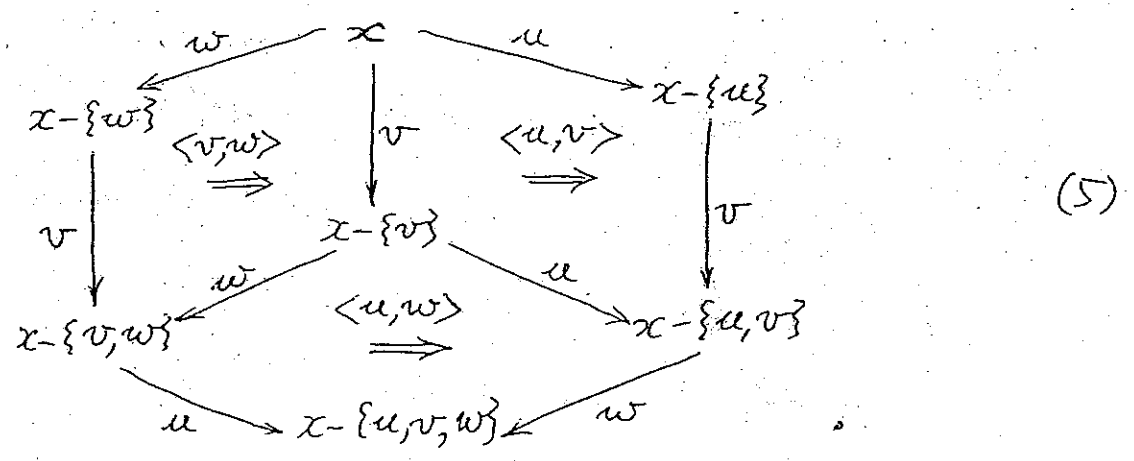
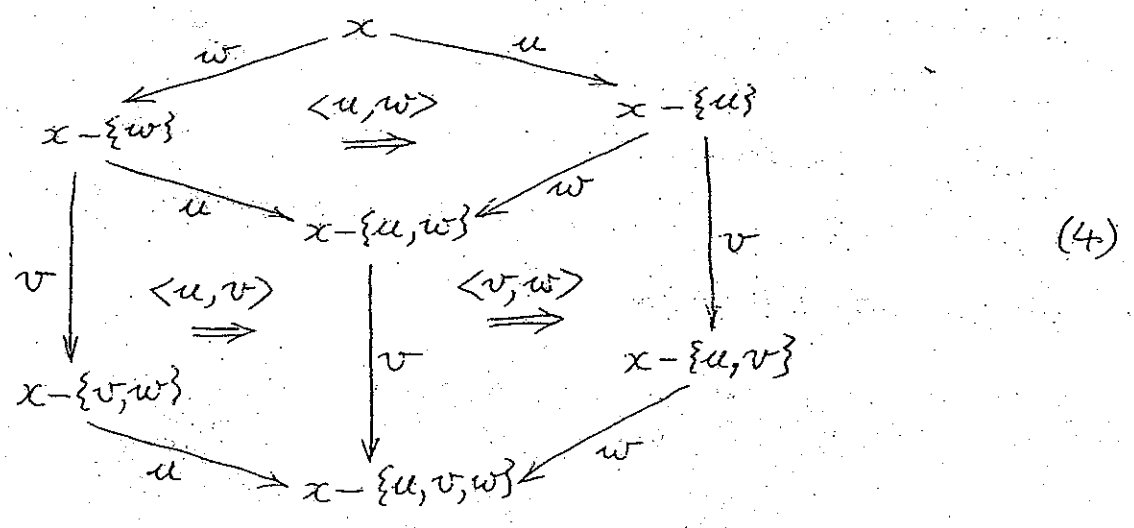
for each  $u, v \in x$  with  $u < v$ . In this situation,

$\sigma(vu) = \mathcal{Q} \subseteq \{(u, v)\} = \sigma(uv)$ ; so we have a morphism of computads  $\mathcal{L}_f(T) \rightarrow \mathcal{Q}(T)$  which extends to a unique 2-functor

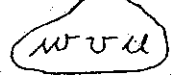
$$h : \mathcal{L}_f(T)^* \rightarrow \mathcal{Q}(T) \tag{3}$$

from the free 2-category on  $\mathcal{L}_f(T)$ .

The following two diagrams denote two particular 2-cells in  $\mathcal{L}_f(T)^*$  for all  $u < v < w$  in  $\mathcal{X}$ :



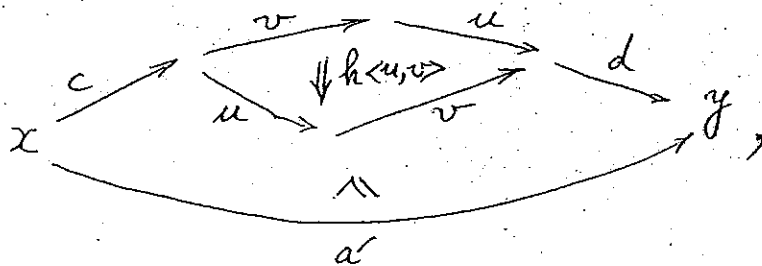
The 2-functor  $h$  takes (4), (5) to the same 2-cell in  $\mathcal{Q}(T)$  since  $\mathcal{Q}(T)$  is locally ordered and the 2-cells (4), (5) have the same source path and same



target path  $uvw$ .

Theorem 1. The 2-functor (3) provides a presentation of  $Q(T)$  as the quotient of the free 2-category  $\mathcal{G}(T)^*$  modulo the relations (4) = (5).

Proof. Consider a non-identity 2-cell  $x \overset{a}{\underset{a'}{\rightrightarrows}} y$  in  $Q(T)$ . Then there exists a consecutive pair  $(v, u)$  in the list  $a$  which occurs in reverse order in the list  $a'$  (or else we would have  $a = a'$ ). If  $v < u$  then  $(v, u) \in \sigma(a) \subseteq \sigma(a')$ , a contradiction. So  $u < v$  and  $(u, v) \in \sigma(a') - \sigma(a)$ . This gives a decomposition of the 2-cell under consideration as



where  $a = cvud$ . By induction on the cardinality of  $\sigma(a') - \sigma(a)$ , the image of  $h$  generates  $Q(T)$ .

It remains to show that the relations (4) = (5) imposed on  $\mathcal{G}(T)^*$  force local order. This is where we need the braid monoid with zero.

A monoid with zero is a monoid  $M$  with an element  $0$  satisfying  $0\mu = \mu 0 = 0$  for all  $\mu \in M$ . A monoid is a set with an associative binary operation having an identity element  $1$ .

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Let  $B_r$  denote the monoid with zero presented by generators  $2, 3, \dots, r$  subject to the relations

$$ji = ij \text{ for } i < j-1 \quad (6)$$

$$(i+1)i(i+1) = i(i+1)i \quad (7)$$

$$ii = 0. \quad (8)$$

Relations (6), (7) are those for Artin's braid group [ ]. So  $B_r$  is called the braid monoid with 0. Recall [ ] that the monoid presented by  $2, \dots, r$  subject to (6), (7) and

$$ii = 1, \quad (8')$$

is the symmetric group  $S_r$  on  $r$  symbols.

For  $i \geq j$ , consider the particular words

$$[i, j] = i(i-1)(i-2)\dots(j+1).$$

They satisfy the equations

$$[i, i] = 1, \quad [i, j][j, k] = [i, k].$$

Eilenberg's result is that each element  $\xi$  of  $B_r$  is either 0 or uniquely of the form

$$\xi = [2, i_2][3, i_3]\dots[r, i_r] \quad (9)$$

where  $1 \leq i_n \leq n$ . It follows that  $B_r$  has cardinality  $r! + 1$ . Notice also that, since the words (9) all represent different permutations (where  $i$  represents the transposition of the  $(i-1)$ -st symbol

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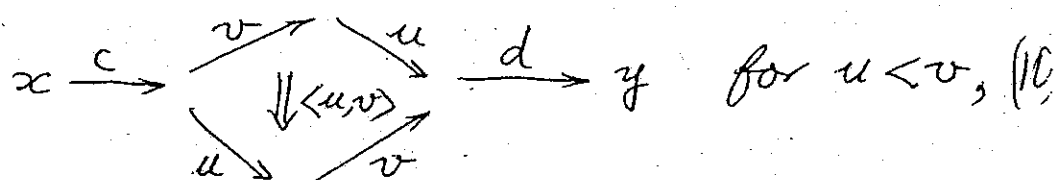
with the  $i$ -th), <sup>Better as action in category of pointed sets.</sup> and since there are  $r!$  of them, every element of  $\mathcal{S}_r$  is uniquely of the form (9).

There is a partial right action of  $\mathcal{B}_r$  on the set of listings  $a$  of  $x-y$  where  $y \subseteq x$  and the cardinality of  $x-y$  is  $r$ . Define  $a * i$  precisely when  $a_{i-1} > a_i$ , and, in that case, to be the list obtained from  $a$  by interchanging  $a_{i-1}$  and  $a_i$ . Extend this associatively to all words in  $2, 3, \dots, r$ ; so we have

$$(a * \xi) * \eta = a * (\xi \eta), \quad a * 1 = a,$$

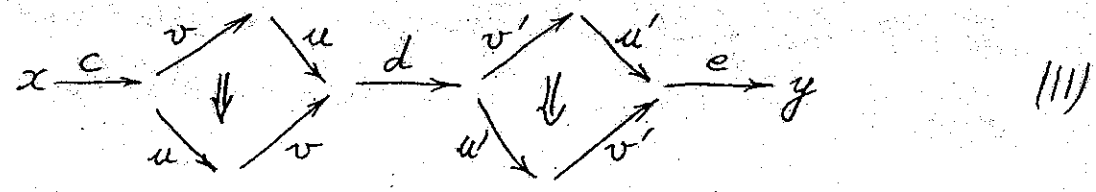
and we declare  $a * 0$  undefined. To see that  $a * \xi = a * \xi'$  when  $\xi, \xi'$  represent the same element of  $\mathcal{B}_r$ , notice that, for each of the relations (6), (7), (8), a list  $a$  can be acted on by the left side if and only if it can be acted on by the right side, and the results are equal when they can be.

The category  $\text{lg}(T)^*(x, y)$  is the quotient of the free category on the graph with edges

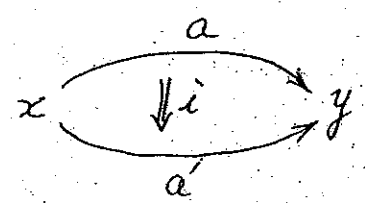


modulo the relations which identify the two paths

from  $cvudv'u'e$  to  $cuvdu'v'e$  obtainable from

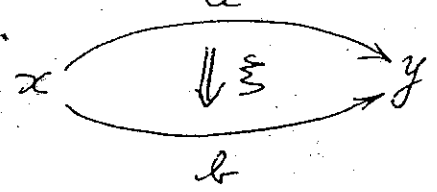


where  $u < v$  and  $u' < v'$ . We can rewrite (10) as



where  $i$  is the length of the list  $cvu$ , where  $a = cvu$  and where  $a * i = a'$ . In this notation, the relations (11) precisely become the relations (6).

Furthermore, the relations (4) = (5) become the relations (7). So, in the quotient of  $\mathcal{L}(T)^*$  by the relations (4) = (5), a 2-cell



is precisely a non-zero element of  $\mathcal{B}_r$  such that  $a * \xi = b$ . But we can assume that  $\xi$  has the normal form (9). By looking at  $b$  we see that  $i_r$  is completely determined. By induction on  $r$ , it follows that  $\xi$  is uniquely determined by  $a$  and  $b$  if it exists. This gives the desired local orderedness.  $\square$

Let  $2\text{-Cat}$  denote the category whose objects are small 2-categories and whose arrows are 2-functors. Let  $\mathcal{C}_u$  denote the full subcategory of  $2\text{-Cat}$  consisting of the objects  $\mathcal{Q}(T)$  for  $T$  a finite totally ordered set.

Theorem 2. The inclusion of  $\mathcal{C}_u$  in  $2\text{-Cat}$  is dense.

Proof. In fact we prove more: we only need those  $\mathcal{Q}(T)$  with the cardinality of  $T$  less than 4. For simplicity we write  $(A, B)$  for the set of 2-functors from  $A$  to  $B$ .

We must show that, for each family of four functions

$$\Theta_n : (\mathcal{Q}(T), A) \longrightarrow (\mathcal{Q}(T), B) \quad 0, 1, 2, 3$$

where  $T$  runs over four totally ordered sets of cardinality  $0 \leq n \leq 3$  such that the square

$$\begin{array}{ccc} (\mathcal{Q}(T), A) & \xrightarrow{\Theta_n} & (\mathcal{Q}(T), B) \\ (f, 1) \downarrow & & \downarrow (f, 1) \\ (\mathcal{Q}(T'), A) & \xrightarrow{\Theta_{n'}} & (\mathcal{Q}(T'), B) \end{array} \quad (12)$$

commutes for all 2-functors  $f: \mathcal{Q}(T') \rightarrow \mathcal{Q}(T)$ , there is a unique 2-functor  $t: A \rightarrow B$  such that  $\Theta = (1, t)$ . So take such a  $\Theta$ .

For  $T$  of cardinality 0 and 1, we can identify  $(\mathcal{Q}(T), A)$  with the set of objects and arrows of  $A$ , respectively. So we define  $t$  on objects and arrows to be given by  $\Theta_0$  and  $\Theta_1$ . By applying (12) with  $f: \mathcal{Q}(0) \rightarrow \mathcal{Q}(1)$



we see that, for  $\alpha: a \rightarrow a'$  an arrow of  $A$ , we have  $t(\alpha): t(a) \rightarrow t(a')$  in  $B$ .

Elements of  $(Q(2), A)$  can be identified with diagrams

$$\begin{array}{ccc}
 a & \xrightarrow{\alpha} & a' \\
 \sigma \downarrow & \xRightarrow{\beta} & \downarrow \sigma' \\
 b & \xrightarrow{\beta} & b'
 \end{array} \tag{13}$$

in  $A$ . By applying (12) with  $f: Q(1) \rightarrow Q(2)$ , we see that  $\Theta_3$  takes (13) to a diagram of the form

$$\begin{array}{ccc}
 t(a) & \xrightarrow{t(\alpha)} & t(a') \\
 t(\sigma) \downarrow & \xRightarrow{\bar{\beta}} & \downarrow t(\sigma') \\
 t(b) & \xrightarrow{t(\beta)} & t(b')
 \end{array}$$

By taking  $f: Q(2) \rightarrow Q(2)$  to be the diagram

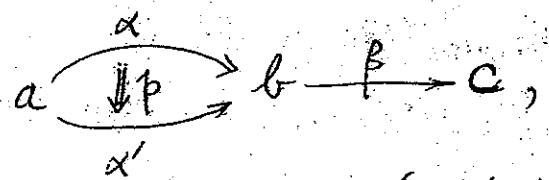
$$\begin{array}{ccc}
 \{u, v\} & \xrightarrow{1} & \{u, v\} \\
 1 \downarrow & \leq & \downarrow uv \\
 \{u, v\} & \xrightarrow{vu} & \mathbb{Q}
 \end{array}$$

in (12) and evaluating at  $s$ , we see that  $\bar{\beta}$  depends only on the 2-cell  $\beta: \sigma\beta \Rightarrow \alpha\sigma'$  so that it can be written  $t(\beta)$ , and that  $t$  preserves composition of arrows.

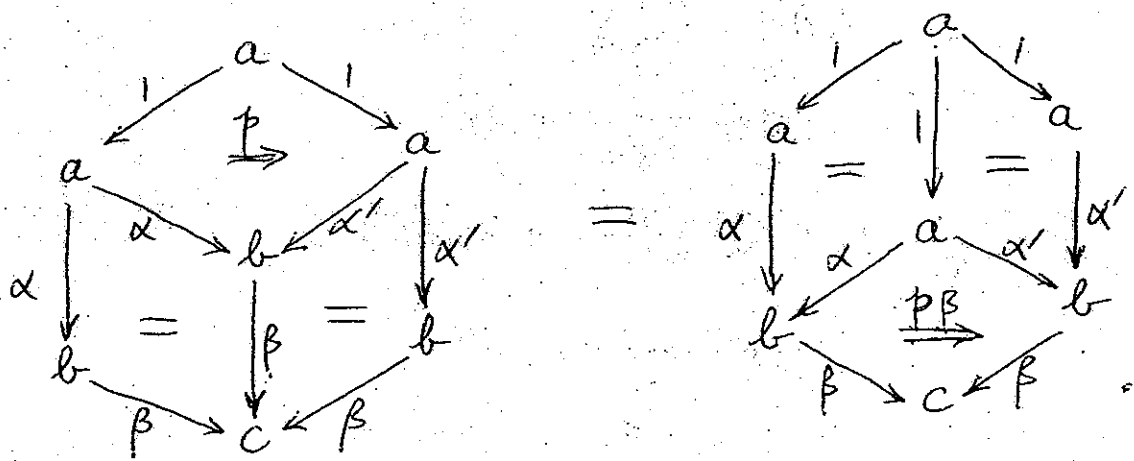
Elements  $h$  of  $(Q(3), A)$  can be identified with

commutative cubes in  $A$  (see (4), (5) with  $x = \{u, v, w\}$ ).  
 By applying (12) with  $f: Q(A) \rightarrow Q(B)$ , we see that the <sup>proper</sup> faces of  $\theta_3(h)$  are just obtained by applying  $t$  to the <sup>proper</sup> faces of  $h$ . So in fact  $\theta_3$  is determined already by  $\theta_0, \theta_1, \theta_2$ ; however, it is needed to prove that  $t$  preserves compositions involving 2-cells.

Given a diagram in  $A$  of the form

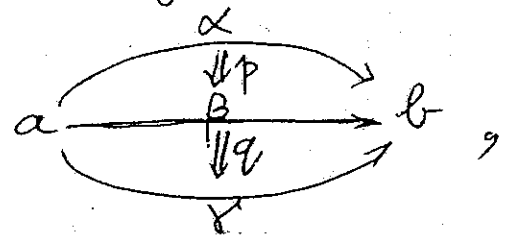


we obtain an element of  $(Q(B), A)$  depicted by

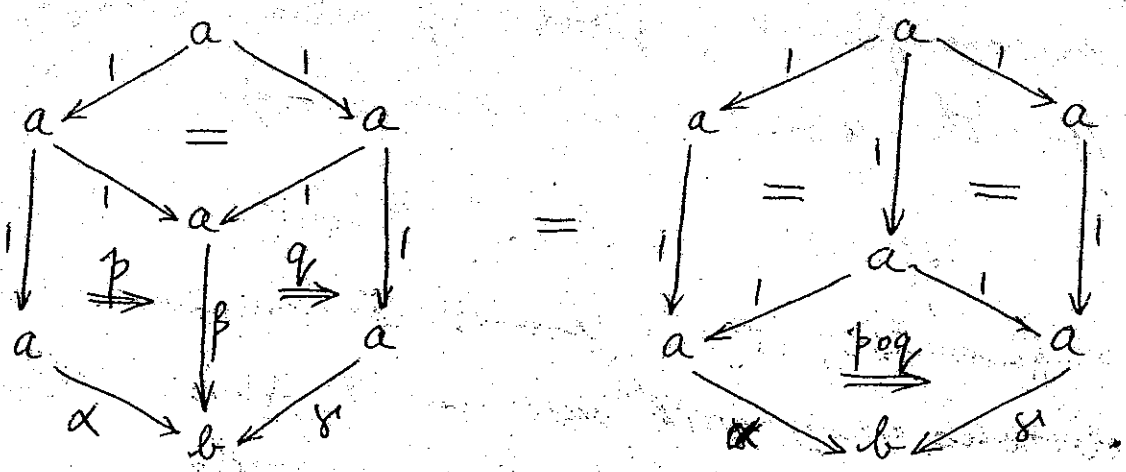


Since  $\theta_3$  takes this to a commutative 3-cube in  $B$ , it follows that  $t(p\beta) = t(p)t(\beta)$ . A symmetric argument applies to show that  $t$  preserves composition of a 2-cell  $p$  with a 1-cell on the left side.

Given a diagram in  $A$  of the form

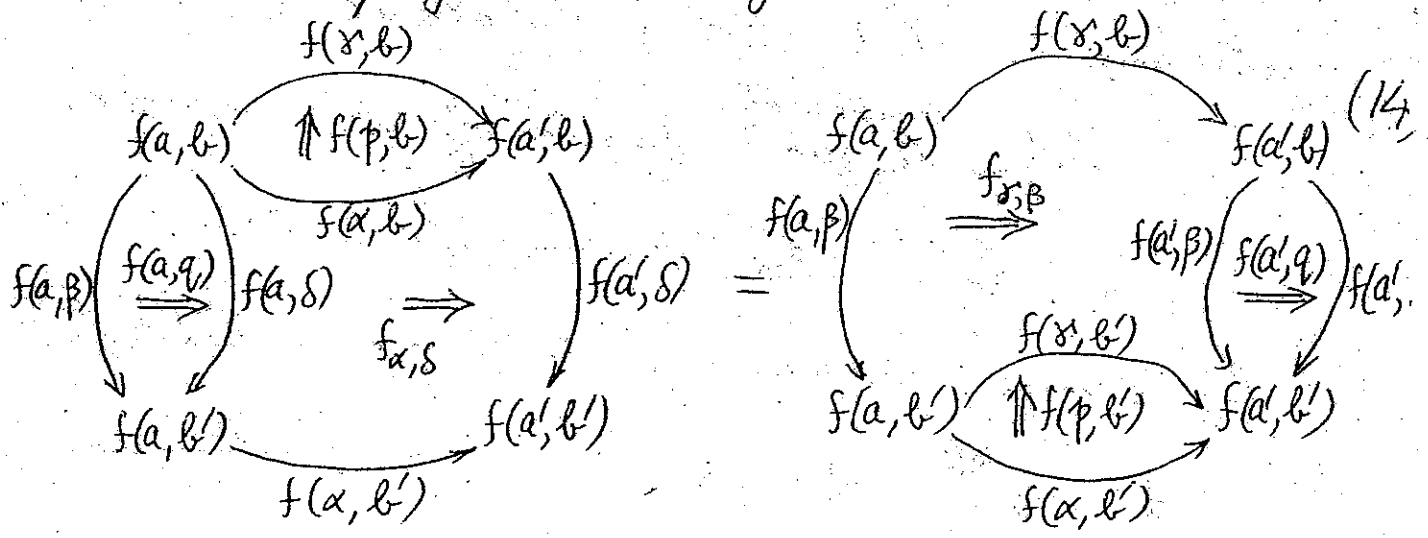


we obtain an element of  $(Q(\mathcal{B}), A)$  of the form



Applying  $\theta_3$  to this, we obtain  $t(p \circ q) = t(p) \circ t(q)$ . So  $\theta$  uniquely determines a 2-functor  $t$ , as required.  $\square$

Recall (Lgray [ , ]) that a quasi-functor  $f: A \times B \rightarrow C$  of two variables where  $A, B, C$  are 2-categories assigns to 2-cells  $a \xrightarrow{\alpha} a'$ ,  $b \xrightarrow{\beta} b'$  in  $A, B$  the data displayed in the diagram



in  $C$  in such a way that equality (14) holds, each  $f(-, b): A \rightarrow C$ ,  $f(a, -): B \rightarrow C$  is a 2-functor, the 2-cell  $f_{\alpha\alpha', \beta}$  is obtained by pasting the square  $f_{\alpha, \beta}$  next to  $f_{\alpha', \beta}$ , the 2-cell  $f_{\alpha, \beta\beta'}$  is

obtained by pasting the square for  $f_{\alpha, \beta}$  on top of the square for  $f_{\alpha, \beta'}$ , and  $f_{\alpha, \beta}$  is an identity 2-cell if either  $\alpha$  or  $\beta$  is an identity arrow.

If  $s: A' \rightarrow A$ ,  $t: B' \rightarrow B$ ,  $r: C \rightarrow C'$  are 2-functors, it is possible to compose in the obvious way with such an  $f: A \times B \rightarrow C$  to obtain a quasi-functor of two variables

$$(s \times t) \circ f \circ r: A' \times B' \rightarrow C'$$

The ordered sum  $S \circ_0 T$  of ordered sets  $S, T$  is the disjoint union of the underlying sets of  $S$  and  $T$  ordered so that it contains  $S, T$  as sub-ordered sets and so that  $u < v$  for all  $u \in S, v \in T$ .

Define the <sup>canonical</sup> quasi-functor of two variables

$$k: Q(S) \times Q(T) \rightarrow Q(S \circ_0 T) \tag{15}$$

by setting, for  $a: x \rightarrow x'$ ,  $b: y \rightarrow y'$  in  $Q(S), Q(T)$ :

$$\begin{array}{ccc}
 k(x, y) & \xrightarrow{k(a, y)} & k(x', y) \\
 k(x, b) \downarrow & \xRightarrow{k_{a, b}} & \downarrow k(x', b) \\
 k(x, y') & \xrightarrow{k(a, y')} & k(x', y')
 \end{array}
 =
 \begin{array}{ccc}
 x + y & \xrightarrow{a} & x' + y \\
 b \downarrow & \leq & \downarrow b \\
 x + y' & \xrightarrow{a} & x' + y'
 \end{array}$$

Theorem 3. Composition with (15) establishes a bijection  
between 2-functors  $g: \mathcal{Q}(S_0; T) \rightarrow A$  and  
quasi-functors of two variables  $f: \mathcal{Q}(S) \times \mathcal{Q}(T) \rightarrow A$

Proof. We begin by examining the computed  $\mathcal{L}(S_0; T)$ .  
 Each object  $z$  can be written uniquely as a disjoint  
 union  $z = x + y$  where  $x \in S$ ,  $y \in T$ . The edges  
 are either of the form  $u: x + y \rightarrow x' + y$  with  
 $u: x \rightarrow x'$  in  $\mathcal{L}(S)$  or of the form  $v: x + y \rightarrow x + y'$   
 where  $v: y \rightarrow y'$  in  $\mathcal{L}(T)$ . The 2-cells are of  
 three kinds (refer to (2)):  $\langle u, u' \rangle$  with  $u < u'$  in  $S$ ,  
 $\langle v, v' \rangle$  with  $v < v'$  in  $T$ , and  $\langle u, v \rangle$  with  $u \in S$ ,  
 $v \in T$ .

Given  $f$ , in order to define  $g$ , we make use  
 of Theorem 1 to see that the equations

$$g(x+y) = f(x, y), \quad g(u) = f(u, y), \quad g(v) = f(x, v),$$

$$g(h\langle u, u' \rangle) = f(\langle u, u' \rangle, y), \quad g(h\langle v, v' \rangle) = f(x, \langle v, v' \rangle),$$

$$\text{and } g(h\langle u, v \rangle) = f_{u, v}$$

determine the unique  $g$  whose composite with  $k$  is  $f$ . [

Using this Theorem 3, we see that 2-functors  
 $s: \mathcal{Q}(S) \rightarrow \mathcal{Q}(S')$ ,  $t: \mathcal{Q}(T) \rightarrow \mathcal{Q}(T')$   
 induce a 2-functor  $s \otimes t$  uniquely determined

by commutativity of the square

$$\begin{array}{ccc}
 \mathcal{Q}(S) \times \mathcal{Q}(T) & \xrightarrow{k} & \mathcal{Q}(S \dot{+} T) \\
 s \times t \downarrow & & \downarrow s \dot{\otimes} t \\
 \mathcal{Q}(S') \times \mathcal{Q}(T') & \xrightarrow{k} & \mathcal{Q}(S' \dot{+} T')
 \end{array} \quad (16)$$

This defines a functor  $\dot{\otimes} : \mathcal{C}u \times \mathcal{C}u \rightarrow \mathcal{C}u$  given by  $\mathcal{Q}(S) \dot{\otimes} \mathcal{Q}(T) = \mathcal{Q}(S \dot{+} T)$  and  $(s, t) \mapsto s \dot{\otimes} t$ .

Let  $\Delta$  denote the category of finite totally ordered sets (including the empty one) with order-preserving functions. We have a functor

$$Q : \Delta \longrightarrow \mathcal{C}u \quad (17)$$

taking  $T$  to  $\mathcal{Q}(T)$ . For an order-preserving  $p : S \rightarrow T$ , the 2-functor  $Q(p) : \mathcal{Q}(S) \rightarrow \mathcal{Q}(T)$  takes  $x$  to  $p(x)$ , takes  $u : x \rightarrow x'$  to  $p(u) : p(x) \rightarrow p(x')$  where  $p(u) \notin p(x')$ , and takes  $\langle u, v \rangle$  as in (2) to  $\langle p(u), p(v) \rangle$  when  $u : x \rightarrow x'$ ,  $v : x' \rightarrow x''$  with  $p(u) \notin p(x')$ ,  $p(v) \notin p(x'')$ .

Ordered sum provides  $\Delta$  with the structure of monoidal category. The canonical

associativity in  $\Delta$  transfers to  $\mathcal{C}_u$  via (17) to yield an isomorphism

$$a : (Q(R) \otimes Q(S)) \otimes Q(T) \cong Q(R) \otimes (Q(S) \otimes Q(T)).$$

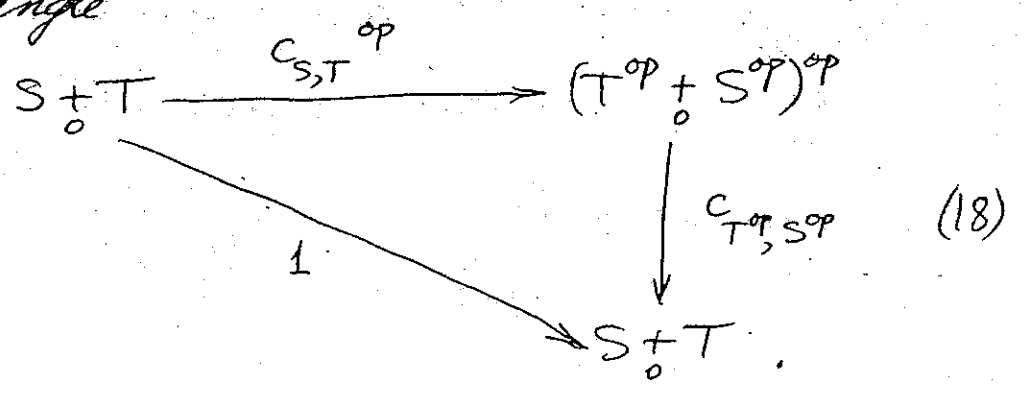
Naturality of  $a$  in  $Q(R), Q(S), Q(T)$  as objects of  $\mathcal{C}_u$  is obvious from (16). All the coherence conditions (including Mac Lane's pentagon) satisfied by the associativity in  $\Delta$  transfer to  $a$  via (17).

Thus we have defined a monoidal structure on the category  $\mathcal{C}_u$  in such a way that (17) becomes a monoidal functor.

The monoidal category  $\Delta$  does not admit a symmetry in the usual sense of Eilenberg-Kelly [ ]; yet there is a canonical isomorphism

$$c : (S \underset{\circ}{+} T)^{op} \cong T^{op} \underset{\circ}{+} S^{op}$$

(where  $S^{op}$  denotes  $S$  with the reverse order) satisfying various coherence conditions including commutativity of the triangle



There is an equally coherent isomorphism

$$c : (Q(S) \otimes Q(T))^{\text{co}} \cong Q(T)^{\text{co}} \otimes Q(S)^{\text{co}} \quad (19)$$

on  $\mathcal{C}_U$  obtained by applying (17), using (16), and using

$$Q(S^{\text{op}}) = Q(S)^{\text{co}}, \quad (20)$$

where  $A^{\text{co}}$  denotes the  $\mathcal{A}$ -category obtained from  $A$  by reversing  $\mathcal{A}$ -cells.

We are now in the familiar situation of a cocomplete category  $\mathcal{A}$ -Cat with a full dense subcategory  $\mathcal{C}_U$  bearing a monoidal structure. It is possible (see Day [ ]) by means of left Kan extension to extend the monoidal structure on  $\mathcal{C}_U$  to one on  $\mathcal{A}$ -Cat. Define the tensor product  $\otimes$  on  $\mathcal{A}$ -Cat to be the  $\mathcal{A}$ -functor in the diagram

$$\begin{array}{ccc} \mathcal{C}_U \times \mathcal{C}_U & \xrightarrow{\quad} & \mathcal{A}\text{-Cat} \times \mathcal{A}\text{-Cat} \\ \otimes \downarrow & \xrightarrow{\cong} & \downarrow \otimes \\ \mathcal{C}_U & \xrightarrow{\quad} & \mathcal{A}\text{-Cat} \end{array}$$

obtained as the left Kan extension of the composite of the left and bottom sides along the top side.

The formula for the tensor product can be written (Day-Kelly [ ], Mac Lane [ ]):



$$A \otimes B = \int^{Q(S), Q(T)} ((Q(S), A) \times (Q(T), B)) \cdot Q(S+T). \quad (21)$$

This means that there is a natural bijection between 2-functors  $A \otimes B \rightarrow C$  and families of functors

$$(Q(S), A) \times (Q(T), B) \longrightarrow (Q(S+T), C) \quad (22)$$

which are natural in  $Q(S), Q(T)$  as objects of  $\mathcal{C}_U$ .

This monoidal structure on 2-Cat is not genuinely symmetric, but we do have a canonical isomorphism

$$c : (A \otimes B)^{co} \cong B^{co} \otimes A^{co} \quad (23)$$

induced by, and as coherent as, (19). We now turn to the question of closedness.

Let  $s, t : A \rightarrow C$  be 2-functors. A lax natural transformation  $\theta : s \rightarrow t$  (also called "quasi-natural" [1]) is given by the data displayed in the square

$$\begin{array}{ccc} s(a) & \xrightarrow{s(\alpha)} & s(a') \\ \theta_a \downarrow & \xRightarrow{\theta_\alpha} & \downarrow \theta_{a'} \\ t(a) & \xrightarrow{t(\alpha)} & t(a') \end{array} \quad (24)$$

satisfying the obvious axioms [1], [2].

Let  $[A, C]$  denote the 2-category whose objects are 2-functors  $s : A \rightarrow C$ , whose arrows are lax natural transformations  $\theta : s \rightarrow t$ , and whose 2-cells

are modifications  $[ \ ]$ ,  $[ \ ]$ ,  $[ \ ]$ .

There is an evaluation quasi-functor of 2-variables

$$e : A \times [A, C] \longrightarrow C \tag{25}$$

given by  $e(a, s) = s(a)$ ,  $e(a, \theta) = \theta_a$ ,  $e(\alpha, s) = s(\alpha)$  and  $e_{\alpha, \theta} = \theta_{\alpha}$  (see (24)). The following result is in Gray [ ]; the proof is just a matter of interpretation and bookkeeping.

Theorem 4. The function taking  $f$  to the composite of  $A \times f$  and  $e$  is a bijection between 2-functors

$f : B \longrightarrow [A, C]$  and quasi-functors of 2-variables  
 $g : A \times B \longrightarrow C$ .  $\square$

It is also clear that quasi-functors of two variables  $g : A \times B \longrightarrow C$  are in bijection with quasi-functors of two variables  $\bar{g} : B^{co} \times A^{co} \longrightarrow C^{co}$  via  $g(a, b) = \bar{g}(b, a)$ .

Theorem 5. There are bijections

$$(B, [A, C]) \cong (A \otimes B, C) \cong (A, [B^{co}, C^{co}]^{co}) \tag{26}$$

natural in  $A, B, C \in 2\text{-Cat}$ .

Proof. Theorem 2 amounts to the formula

$$A \cong \int^{Q(S)} (Q(S), A) \cdot Q(S). \tag{27}$$

Together Theorems 3 and 4 yield

$$(Q(T), [Q(S), C]) \cong (Q(S) \otimes Q(T), C). \quad (28)$$

Theorem 4 and the remark following it yield

$$(B, [A, C]) \cong (A, [B^{co}, C^{co}])^{co}. \quad (29)$$

The calculation

$$\begin{aligned} (B, [Q(S), C]) &\stackrel{(27)}{\cong} \int^{Q(T)} (Q(T), B) \cdot Q(T), [Q(S), C] \\ &\cong \int_{Q(T)} (Q(T), [Q(S), C])^{(Q(T), B)} \quad (\text{by definition of coend}) \\ &\stackrel{(28)}{\cong} \int_{Q(T)} (Q(S) \otimes Q(T), C)^{(Q(T), B)} \\ &\cong \int^{Q(T)} (Q(T), B) \cdot (Q(S) \otimes Q(T), C) \quad (\text{by definition of coend}) \\ &\cong (Q(S) \otimes B, C) \quad (\text{by (21) and Yoneda}). \end{aligned}$$

Thus we have (26) for  $A = Q(S)$ .

A similar calculation using (29) gives the result.  $\square$

Thus we have described a biclosed monoidal structure on the category  $2\text{-Cat}$ . There is another one which we shall now describe.

Let  $\{A, C\}$  denote the sub-2-category of  $[A, C]$  consisting of the same objects, the lax natural

transformations  $\theta: s \rightarrow t$  for which each  $\theta_\alpha$  (see (24)) is invertible (these  $\theta$  are called pseudo-natural transformations), and all modifications. There is an isomorphism of 2-categories

$$\{A, C\} \cong \{A^{co}, C^{co}\}^{co} \quad (30)$$

obtained by replacing  $\theta_\alpha$  by  $\theta_\alpha^{-1}$ .

The tensor product  $A \boxtimes B$  for this smaller internal hom will now be constructed. The non-identity 2-cell of the "square"  $Q(\mathcal{Q})$  can be regarded as a modification

$$1 \begin{array}{c} \xrightarrow{\quad} \\ \Downarrow \Rightarrow \Downarrow \\ \xrightarrow{\quad} \end{array} Q(\mathcal{Q}). \quad (31)$$

Arrows  $\alpha: a \rightarrow a'$ ,  $\beta: b \rightarrow b'$  in  $A, B$  can be regarded as 2-functors  $\bar{\alpha}: Q(1) \rightarrow A$ ,  $\bar{\beta}: Q(1) \rightarrow B$ , so, we obtain a 2-functor

$$Q(\mathcal{Q}) \cong Q(1) \otimes Q(1) \xrightarrow{\bar{\alpha} \otimes \bar{\beta}} A \otimes B. \quad (32)$$

Combining (31), (32), we obtain a modification

$$\text{arr } A \times \text{arr } B \begin{array}{c} \xrightarrow{\quad} \\ \Downarrow \Rightarrow \Downarrow \\ \xrightarrow{\quad} \end{array} A \otimes B, \quad (33)$$

where  $\text{arr } A$  denotes the set of arrows of  $A$ . Let

$$A \otimes B \longrightarrow A \boxtimes B \quad (34)$$

be the universal 2-functor which inverts the

modification (33).

Theorem 6. The first bijection of Theorem 5 induces a bijection

$$(B, \{A, C\}) \cong (A \boxtimes B, C)$$

natural in  $A, B, C \in \mathcal{A}\text{-Cat}$ .

Proof. Write  $[A, C]$  for the sub- $\mathcal{A}$ -category of  $\llbracket A, C \rrbracket$  where we restrict to arrows which are  $\mathcal{A}$ -natural transformations (each  $\theta_\alpha$  an identity). It is easy to see that

$$[B, \{A, C\}] \longrightarrow [B, \llbracket A, C \rrbracket]$$

is the universal  $\mathcal{A}$ -functor which inverts the modification

$$[B, \llbracket A, C \rrbracket] \xrightarrow{\quad \Downarrow \Rightarrow \Downarrow \quad} [A \boxtimes B, C].$$

This gives a  $\mathcal{A}$ -natural isomorphism of  $\mathcal{A}$ -categories

$$[B, \{A, C\}] \cong [A \boxtimes B, C];$$

the object bijection gives the result.  $\square$

Theorem 7. The tensor product  $\boxtimes$  (together with associativity isomorphism induced from  $\otimes$  and symmetry induced from (23) along (34)) enriches  $\mathcal{A}\text{-Cat}$  with a symmetric monoidal closed structure.

Proof. It is easy to see that the associativity isomorphism for  $\otimes$  induces one for  $\boxtimes$ ; coherence

follows since (34) is epic. Combining (30), (18), (23), we obtain all that is needed for a symmetry except for commutativity of the triangle (omitting associativities)

$$\begin{array}{ccc}
 A \boxtimes B \boxtimes C & \xrightarrow{c} & C \boxtimes A \boxtimes B \\
 \searrow^{1 \boxtimes c} & & \nearrow_{c \boxtimes 1} \\
 & A \boxtimes C \boxtimes B &
 \end{array}$$

which could not even be drawn in the  $\otimes$  case with  $( )^{co}$  involved. Now that it can be drawn, it is easy. [