# Making weak maps compose strictly

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#### **Outline**

**Motivation** 

Weak maps of bicategories

Weak maps of tricategories

Weak maps of weak  $\omega$ -categories

(NB: Talk notes available at http://www.dpmms.cam.ac.uk/ $\sim$ rhgg2)

### **Motivation**

### Consider a category & with:

- Objects being higher-dimensional ——s;
- ► *Morphisms* being strict structure-preserving maps.

### Would like to derive $\mathscr{C}_{wk}$ with:

- ► Same objects;
- ► Morphisms being *weak* structure-preserving maps.

### **Motivation**

Idea from homotopy theory: identify

weak maps 
$$X \to Y$$
 with strict maps  $X' \to \tilde{Y}$ 

#### where:

- ightharpoonup X' is a cofibrant replacement for X;
- $\tilde{Y}$  is a fibrant replacement for Y.

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- ▶ A strict map  $X \to \tilde{Y}$  is a map which preserves the differential only up to homotopy;

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- ▶ A strict map  $X \to \tilde{Y}$  is a map which preserves the differential only up to homotopy;
- ▶ A strict map  $X' \to \tilde{Y}$  (= weak map  $X \to Y$ ) is a map which preserves the R-module structure and the differential only up to homotopy.

Let  $f: X \to Z$  in **Cat**. Can form the *interval category*  $f/\mathbf{Cat}/Z$ :

- ▶ *Objects* are  $X \xrightarrow{g} Y \xrightarrow{h} Z$  with hg = f;
- ► *Morphisms* are commutative diamonds:



Corresponding weak maps should be pseudo-commutative diamonds:



Corresponding weak maps should be pseudo-commutative diamonds:



... and these are precisely strict maps



where:

... fibrant replacement of  $X \xrightarrow{j} W \xrightarrow{k} Z$  is:

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and cofibrant replacement of  $X \xrightarrow{g} Y \xrightarrow{h} Z$  is:

$$X \xrightarrow{l_g} g \uparrow_{\cong} Y \xrightarrow{h \circ r_g} Z.$$

### How to compose weak maps?

### Idea from category theory:

- ▶ Cofibrant replacement should be a *comonad* (–)':  $\mathscr{C} \to \mathscr{C}$ ;
- ▶ Fibrant replacement should be a monad (-):  $\mathscr{C} \to \mathscr{C}$ ;
- ▶ There should be a distributive law  $d_X: (\tilde{X})' \to \widetilde{X}'$ .

## How to compose weak maps?

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- ▶ Fibrant replacement should be a *monad*  $\widetilde{(-)}$ :  $\mathscr{C} \to \mathscr{C}$ ;
- ▶ There should be a distributive law  $d_X$ :  $(\tilde{X})' \to \tilde{X}'$ .

Now composition of weak maps is two-sided Kleisli composition:

$$\begin{array}{ccc} (X' \xrightarrow{f} \tilde{Y}) \ \circ \ (Y' \xrightarrow{g} \tilde{Z}) & := \\ X' \xrightarrow{\Delta_X} X'' \xrightarrow{f'} (\tilde{Y})' \xrightarrow{d_Y} \widetilde{Y'} \xrightarrow{\tilde{g}} \tilde{Z} \xrightarrow{\mu_Z} \tilde{Z}. \end{array}$$

- ► Cofibrant replacement is a comonad [Grandis—Tholen 2006];
- ► Fibrant replacement *is* a monad [loc. cit.];
- ► There *is* a distributive law between them;

and corresponding Kleisli composition is what you think it is: pasting of pseudo-commutative diamonds.

### In general

If a (locally presentable) category  ${\mathscr C}$  has a cofibrantly generated model structure on it, then:

- Cofibrant replacement can be made a comonad [G. 2008];
- ► Fibrant replacement can be made a monad [loc. cit.];
- ▶ But not clear how to get a distributive law between them!

So in this talk, we focus on the case where *every object is fibrant*. (As then we only need cofibrant replacement comonad).

## Weak maps of bicategories

Consider the category **Bicat**<sub>s</sub>:

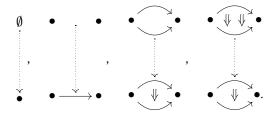
- Objects are bicategories;
- ► *Morphisms* are strict homomorphisms.

There is a cofibrantly generated model structure on **Bicat**<sub>s</sub> [Lack, 2004], wherein:

- Weak equivalences are biequivalences;
- ► Every object is fibrant.

What are the corresponding weak maps?

First we describe cofibrant replacement comonad (–)'. It is generated by the following set of maps in **Bicat**<sub>s</sub>:



### Explicitly, if $\mathfrak{B}$ is a bicategory, then $\mathfrak{B}'$ is given as follows:

▶ Ignore the 2-cells and form the free bicategory  $FU\mathcal{R}$  on the underlying 1-graph of  $\mathcal{R}$ ;

Explicitly, if  $\mathcal{B}$  is a bicategory, then  $\mathcal{B}'$  is given as follows:

- ▶ Ignore the 2-cells and form the free bicategory *FU*ℬ on the underlying 1-graph of ℬ;
- ▶ Factorise the counit map  $\varepsilon$ :  $FU\Re \to \Re$  as

$$FU\mathfrak{B} \xrightarrow{a} \mathfrak{B}' \xrightarrow{b} \mathfrak{B}$$

where a is bijective on objects and 1-cells and b is locally fully faithful.

(NB: this is the flexible replacement of [Blackwell-Kelly-Power 1989]).

### **Proposition (Coherence for homomorphisms)**

The co-Kleisli category of (-)': **Bicat**<sub>s</sub>  $\rightarrow$  **Bicat**<sub>s</sub> is isomorphic to the category **Bicat** of bicategories and homomorphisms.

#### **Proposition (Coherence for homomorphisms)**

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#### Proof.

- ▶ First define a comonad H on  $\mathbf{Bicat}_s$  such that  $\mathbf{Kl}(H) \cong \mathbf{Bicat}$  by construction;
- ▶ Then show that  $H \cong (-)'$  as comonads.

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### Proof.

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- ▶ Then show that  $H \cong (-)'$  as comonads.

Explicitly, given bicategory  $\Re$ , we form  $H\Re$  as follows:

Start with  $FU\mathcal{B}$  as above. Given  $f: X \to Y$  in  $\mathcal{B}$ , write  $[f]: X \to Y$  for corresponding generator in  $FU\mathcal{B}$ .

Now adjoin 2-cells to  $FU\Re$  as follows:

► For each 
$$\alpha$$
:  $f \Rightarrow g$  in  $\Re$ , a 2-cell

$$[\alpha]: [f] \Rightarrow [g];$$

► For each 
$$X \in \mathcal{B}$$
, a 2-cell  $\eta_X : id_X \Rightarrow [id_X]$ ;

► For each 
$$X \xrightarrow{f} Y \xrightarrow{g} Z \in \mathcal{B}$$
, a 2-cell
$$\mu_{\sigma, f} \colon [g] \circ [f] \Rightarrow [g \circ f];$$

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$$\mu_{g,f} \colon [g] \circ [f] \Rightarrow [g \circ f];$$

And quotient out the 2-cells by equations making:

- ► [–] be functorial on 2-cells;
- $\blacktriangleright \mu_{g,f}$  be natural in g and f;
- ► The  $\mu_{g,f}$ 's and  $\eta_X$ 's satisfy the unit and associativity laws.

The result of this is H $\Re$ .

▶ By construction, maps H $\Re$  →  $\mathscr{C}$  are in bijection with

homomorphisms  $\mathfrak{B} \to \mathscr{C}$ .

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(comonad structure on *H* is combinatorial essence of composition

of homomorphisms—compare [Hess-Parent-Scott 2006]).

- ▶ By construction, maps H $\Re$  →  $\mathscr{C}$  are in bijection with
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- ► We can now make H into a comonad so that  $\mathbf{Kl}(H) \cong \mathbf{Bicat}$

(comonad structure on *H* is combinatorial essence of composition

▶ Finally, we show that  $H \cong (-)'$  as comonads (a normalization proof).

of homomorphisms—compare [Hess-Parent-Scott 2006]).

## Weak maps of tricategories

Consider the category **Tricat**<sub>s</sub>:

- ► *Objects* are tricategories;
- ► *Morphisms* are strict homomorphisms.

We use an algebraic definition of tricategory, so  $\mathbf{Tricat}_s$  is l.f.p. and in particular cocomplete.

## Weak maps of tricategories

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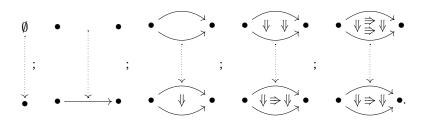
We use an algebraic definition of tricategory, so  $\mathbf{Tricat}_s$  is l.f.p. and in particular cocomplete.

No-one has written down the cofibrantly generated model structure on  $Tricat_s$  yet, but it should have:

- Weak equivalences being triequivalences;
- Every object being fibrant.

Can we describe the corresponding weak maps?

Yes: because we can describe the cofibrant replacement comonad (-)'. It's generated by the following set of maps in **Tricat**<sub>3</sub>:



Length the 2 and 2 calls and form the free trigategory FIIT on the

Explicitly, if  $\mathcal{T}$  is a tricategory, then  $\mathcal{T}'$  is given as follows:

▶ Ignore the 2- and 3-cells and form the free tricategory  $FU\mathcal{T}$  on the underlying 1-graph of  $\mathcal{T}$ . Write  $\varepsilon \colon FU\mathcal{T} \to \mathcal{T}$  for the counit map.

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- underlying 1-graph of  $\mathcal{T}$ . Write  $\varepsilon \colon FU\mathcal{T} \to \mathcal{T}$  for the counit map.
- ▶ For each pair of 1-cells  $f, g: X \to Y$  in  $FU\mathcal{T}$  and each 2-cell  $\alpha : \varepsilon(f) \Rightarrow \varepsilon(g)$  in  $\mathcal{T}$ , adjoin a 2-cell  $(f, g, \alpha) : f \Rightarrow g$  to  $FU\mathcal{T}$ . Call

Explicitly, if  $\mathcal{T}$  is a tricategory, then  $\mathcal{T}'$  is given as follows:

- ▶ Ignore the 2- and 3-cells and form the free tricategory  $FU\mathcal{T}$  on the
  - underlying 1-graph of  $\mathcal{T}$ . Write  $\varepsilon \colon FU\mathcal{T} \to \mathcal{T}$  for the counit map. ▶ For each pair of 1-cells  $f, g: X \to Y$  in  $FU\mathcal{T}$  and each 2-cell
  - $\alpha : \varepsilon(f) \Rightarrow \varepsilon(g)$  in  $\mathcal{T}$ , adjoin a 2-cell  $(f, g, \alpha) : f \Rightarrow g$  to  $FU\mathcal{T}$ . Call the result  $\mathcal{T}^{\#}$ , and write  $\varepsilon^{\#}: \mathcal{T}^{\#} \to \mathcal{T}$  for the induced counit.
  - Factorise  $\varepsilon^{\#}$  as  $\mathfrak{I}^{\#} \xrightarrow{a} \mathfrak{I}' \xrightarrow{b} \mathfrak{I}$

where a is bijective on 0-, 1- and 2-cells and b is locally locally fully

faithful.

► As before (–)′ underlies a comonad, so we obtain a category **Kl**<sub>(−)′</sub> of "tricategories and weak maps".

- ► As before (–)' underlies a comonad, so we obtain a category **Kl**<sub>(−)'</sub> of "tricategories and weak maps".
- ► *A priori* quite surprising, because trihomomorphisms à la [Gordon-Power-Street 1995] do not compose associatively: there *is* no category of tricategories and (ordinary) trihomomorphisms.

- ► As before (–)′ underlies a comonad, so we obtain a category **Kl**<sub>(−)′</sub> of "tricategories and weak maps".
- ► *A priori* quite surprising, because trihomomorphisms à la [Gordon-Power-Street 1995] do not compose associatively: there *is*
- no category of tricategories and (ordinary) trihomomorphisms.So what *do* these new weak morphisms look like? Can they really

be as weak as trihomomorphisms?

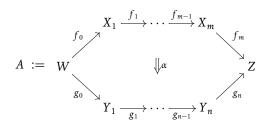
#### **Definition**

An unbiased trihomomorphism  $F: \mathcal{T} \to \mathcal{U}$  is given by:

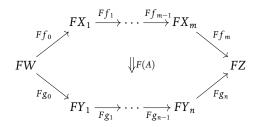
- For each object  $X \in \mathcal{T}$ , an object  $FX \in \mathcal{U}$ ;
- ▶ For each 1-cell  $f: X \to Y \in \mathcal{T}$ , a 1-cell  $Ff: FX \to FY$  in  $\mathcal{U}$ ;

Plus...

► For every bracketed pasting diagram

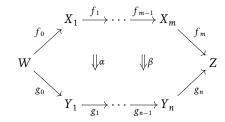


in  $\mathcal{T}$ , a 2-cell



 $\text{in } \mathcal{U}.$ 

▶ For every pair of bracketed pasting diagrams A, B with the same boundary in  $\mathcal{T}$ ; i.e.,

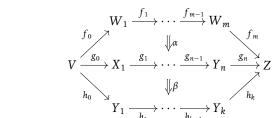


and for every 3-cell  $\Gamma$ :  $\alpha \Rightarrow \beta$  between them, a 3-cell

$$F\Gamma \colon F(A) \Rightarrow F(B)$$

in  $\mathcal{U}$ .

► For every pair of composable bracketed pasting diagrams *A*, *B* in *T*; i.e.,



a 3-cell

$$\mu_{A,B} \colon F(B) \cdot F(A) \Rightarrow F(B \cdot A)$$

in  ${\mathscr U}.$ 

▶ For every identity pasting diagram  $id_{\{f_i\}}$ , i.e.

$$Y_1 \xrightarrow{f_1} \cdots \xrightarrow{f_{m-1}} Y_m$$

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$$X \qquad \qquad \downarrow \text{id} \qquad Z$$

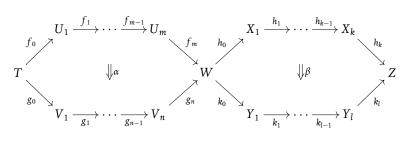
$$Y_{1} \xrightarrow{f_{1}} \cdots \xrightarrow{f_{m-1}} Y_{m}$$

 $\eta_{\{f_i\}} : \mathrm{id}_{\{F_f\}} \Rightarrow F(\mathrm{id}_{\{f_i\}})$ 

in  $\mathcal{T}$ , a 3-cell

in  $\mathcal{U}$ .

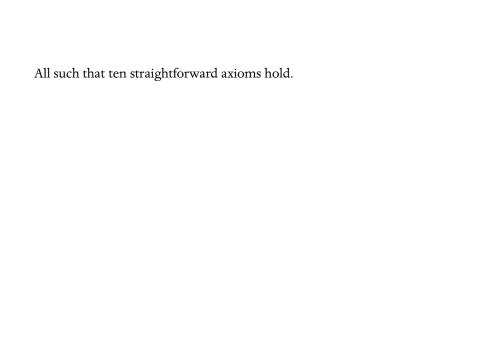
► For every pair *A*, *B* of horizontally composable bracketed pasting diagrams:



in  $\mathcal{T}$ , a 3-cell

$$\gamma_{A,B} \colon F(B) \otimes F(A) \Rightarrow F(B \otimes A)$$

in  $\mathcal{U}$ .



All such that ten straightforward axioms hold.

#### Proposition (Coherence for unbiased trihomomorphisms)

The co-Kleisli category of (-)': **Tricat**<sub>s</sub>  $\rightarrow$  **Tricat**<sub>s</sub> is isomorphic to the category **UTricat** of tricategories and unbiased trihomomorphisms.

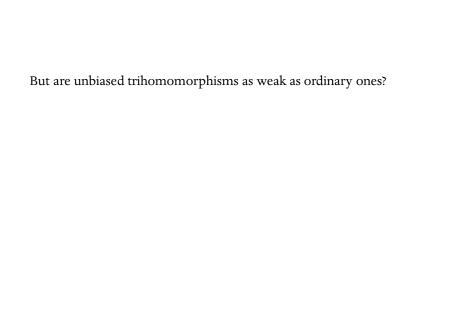
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## Proof.

- ▶ First define a comonad H on  $Tricat_s$  such that  $Kl(H) \cong UTricat$  by construction;
- ▶ Then show that  $H \cong (-)'$  as comonads.



But are unbiased trihomomorphisms as weak as ordinary ones?

#### **Proposition**

Every unbiased trihomomorphism  $\mathcal{T} \to \mathfrak{U}$  gives rise to an ordinary trihomomorphism; and vice versa.

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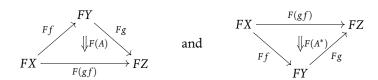
Every unbiased trihomomorphism  $\mathcal{T} \to \mathfrak{U}$  gives rise to an ordinary trihomomorphism; and vice versa.

E.g., given an unbiased trihomomorphism  $F \colon \mathcal{T} \to \mathcal{U}$  let us show that it preserves 1-cell composition up to equivalence.

Let  $X \xrightarrow{f} Y \xrightarrow{g} Z$  in  $\mathcal{T}$ . We have bracketed pasting diagrams  $A, A^* :=$ 

$$X \xrightarrow{gf} Z$$
 and  $X \xrightarrow{gf} Z$   $X \xrightarrow{gf} Z$ 

in  $\mathcal{T}$ , and so obtain 2-cells



in  $\mathcal{U}$ . Moreover,  $A \circ A^* = \text{id implies } F(A) \circ F(A)^* \cong \text{id and dually.}$ 

Can formalise the above equivalence using a <i>bicategory of tricategories</i> .

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## Definition (G.-Gurski 2008)

Given ordinary trihomomorphisms  $F,G: \mathcal{T} \to \mathcal{U}$ , a tricategorical icon  $\Gamma: F \Rightarrow G$ :

- ► Exists only if *F* and *G* agree on 0- and 1-cells;
- ▶ Is then given by 3-cells  $\Gamma_{\alpha}$ :  $F\alpha \Rightarrow G\alpha$  for each 2-cell  $\alpha \in \mathcal{T}$ ;
- ▶ Plus some coherence data.

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- ▶ Plus some coherence data.

Similar definition of unbiased tricategorical icon.

There is a bicategory **Tricat**<sub>2</sub> with:

- ► Objects being tricategories;
- ► Morphisms being (ordinary) trihomomorphisms;
- ▶ 2-cells being tricategorical icons.

There is a bicategory **Tricat**<sup>2</sup> with:

- Objects being tricategories;
- Morphisms being (ordinary) trihomomorphisms;
- 2-cells being tricategorical icons.

There is also a 2-category **UTricat**<sub>2</sub> with:

- ► Objects being tricategories;
- Morphisms being unbiased trihomomorphisms;
- 2-cells being unbiased tricategorical icons.

The bicategory  $Tricat_2$  is equivalent to the 2-category  $UTricat_2$ .

The bicategory  $Tricat_2$  is equivalent to the 2-category  $UTricat_2$ .

## Proof.

- ► First extend the category **Tricat**<sub>s</sub> to a 2-category, with icons as 2-cells;
- ► Then define a 2-comonad H on  $\mathbf{Tricat}_s$  such that  $\mathbf{Kl}(H) \cong \mathbf{UTricat}_2$  by construction;
- ► Then define a pseudo-comonad K on  $\mathbf{Tricat}_s$  such that  $\mathbf{Kl}(K) \cong \mathbf{Tricat}_2$  by construction;
- ▶ Finally show that  $H \simeq K$  as pseudo-comonads.

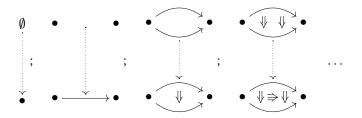
# Weak maps of weak $\omega$ -categories

Consider the category  $\omega$ -**Cat**<sub>s</sub>:

- *Objects* are (algebraic) weak  $\omega$ -categories;
- ► *Morphisms* are strict homomorphisms.

We can play the same game as before to obtain a category  $\omega$ -**Cat** of weak  $\omega$ -categories and weak homomorphisms.

This time the cofibrant replacement comonad (–)' is generated by the following set of maps in  $\omega$ -Cat<sub>s</sub>:



Explicitly, (–)':  $\omega$ -Cat<sub>s</sub>  $\rightarrow \omega$ -Cat<sub>s</sub> is the comonad arising from the adjunction

$$\omega$$
-Cptd  $\xrightarrow{U}_{E} \omega$ -Cat<sub>s</sub>

where  $\omega$ -**Ctpd**<sub>s</sub> is the category of  $\omega$ -computads.

And by now natural to define...

Explicitly, (-)':  $\omega$ -**Cat**<sub>s</sub>  $\to \omega$ -**Cat**<sub>s</sub> is the comonad arising from the adjunction

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where  $\omega$ -**Ctpd**, is the category of  $\omega$ -computads.

And by now natural to define...

#### **Definition**

The category  $\omega$ -**Cat** of weak  $\omega$ -categories and weak morphisms is the co-Kleisli category of (–)':  $\omega$ -**Cat**<sub>s</sub>  $\rightarrow \omega$ -**Cat**<sub>s</sub>.

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